
Expressions and Assignment statements

CS 315 – Programming Languages

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Introduction

- Expressions are the fundamental means of specifying computations in a programming language
- Syntax of Expressions – BNFs
- Semantic of Expressions – will be discussed in this chapter
- To understand expression evaluation, need to be familiar with the orders of operator and operand evaluation
- Essence of imperative languages is the dominant role of assignment statements

Arithmetic expressions

- Arithmetic evaluation was one of the motivations for the development of the first programming languages
- Arithmetic expressions consist of operators, operands, parentheses, and function calls

Arithmetic expressions : Design Issues

- operator precedence rules
- operator associativity rules
- order of operand evaluation
- operand evaluation side effects
- operator overloading
- mode mixing expressions

Arithmetic expressions: Operators

- A **unary** operator has one operand
- A **binary** operator has two operands
- A **ternary** operator has three operands

Arithmetic expressions: Operator Precedence Rules

- The *operator precedence rules* for expression evaluation define the order in which “adjacent” operators of different precedence levels are evaluated
 - $3 + 4 * 5$ (35 or 23)
- Typical precedence levels
 - parentheses
 - unary operators
 - ****** (if the language supports it)
 - *****, **/**
 - **+**, **-**
- Usually Unary minus (-) should not be adjacent to another operator.
 - Example, $A+-B*C$ is usually illegal. It is legal in C.
- APL has a single level of precedence rules

Arithmetic Expressions: Operator Associativity Rule

- The *operator associativity rules* for expression evaluation define the order in which adjacent operators with the same precedence level are evaluated
- $A - B + C - D$
- if all $+$ and $-$ are at the **same precedence level** then what is the order of evaluation?
- Typical associativity rule – Left to right
- Exceptions
 - In **FORTRAN**, exponentiation ($**$) is right associative
 - In Ada, exponentiation ($**$) is nonassociative; $A**B**C$ is illegal.
 - In **C**: prefix $++$, prefix $--$, unary $+$, unary $-$ and $=$ are right associative.
 - Sometimes unary operators associate right to left (e.g., in FORTRAN)
- APL is different; all operators have equal precedence and all operators associate right to left
 - $A * B + C$ is evaluated as $A * (B + C)$
- Precedence and associativity rules can be overridden with parentheses

Arithmetic Expressions: Parantheses

- Precedence and associativity rules can be altered by placing parantheses
- Example: $(A+B)*C$

Arithmetic Expressions: Conditional Expressions

- Conditional Expressions
 - C-based languages (e.g., C, C++)
 - An example:

```
average = (count == 0)? 0 : sum / count
```

- Evaluates as if written like

```
if (count == 0) average = 0  
else average = sum / count
```

Arithmetic Expressions: Operand Evaluation Order

- Variables: fetch the value from memory
 - Constants: sometimes a fetch from memory; sometimes the constant is in the machine language instruction
 - Parenthesized expressions: evaluate all operands and operators first
-
- If the operands **do not have side effects** then the operand evaluation order **does not matter**.

Arithmetic Expressions: Potentials for Side Effects

- *Functional side effects*: when a function changes a two-way parameter or a non-local variable
- Problem with functional side effects:
 - When a function referenced in an expression alters another operand of the expression; e.g., for a parameter change:
 `a = 10;`

 /* assume that fun changes its parameter */
 `b = a + fun(a);`

Arithmetic Expressions: Side Effects

- Side effects of a function call (**functional side effect**):
- Function changes either one of its parameters or a global variable.

Arithmetic Expressions: Side Effects

- Example in PASCAL:
- `function foo (var x: real): real;`
- `begin`
- `x := x/2; /* similar situation occurs when */`
- `foo := x; /* function changes a global var */`
- `end;`
- ...
- `A := 10;`
- `B := A + foo(A);`
- ...
- If A is fetched first, then `foo(A)` is evaluated, the result is 15
- If `foo(A)` is evaluated first, then the results is 10

Arithmetic Expressions: Side Effects

- Example in C:
- `int a = 5;`
- `int foo(){`
- `a=7;`
- `return 3;`
- `} /* foo */`
- `main () {`
- `a = a + foo();`
- `printf("a: %d\n", a);`
- `}`

- When compiled with gcc prints **a:10**
- **When compiled with cc prints a:8**

Functional Side Effects

- Two possible solutions to the problem
 1. Write the language definition to disallow functional side effects
 - No two-way parameters in functions
 - No non-local references in functions
 - **Advantage:** it works!
 - **Disadvantage:** inflexibility of two-way parameters and non-local references
 2. Write the language definition to demand that operand evaluation order be fixed (E.g., Java: left to right)
 - **Disadvantage:** limits some compiler optimizations

Overloaded Operators

- Use of an operator for more than one purpose is called *operator overloading*
- Some are common (e.g., `+` for `int` and `float`)
- Some are potential trouble (e.g., `*` in C and C++)
 - Loss of compiler error detection (omission of an operand should be a detectable error)
 - Some loss of readability
 - Can be avoided by introduction of new symbols (e.g., Pascal's **`div`** for integer division)

Overloaded Operators

- In **C**: **&** as a binary operator: bitwise logical AND
- as a unary operator: address of a variable.
- Two unrelated meanings. **Not readable.**
- Example,
- $x = z \ \& \ y$
- If the programmer forgets to type z it is
- $x = \& \ y$
- Compiler cannot detect such error.

Overloaded Operators

- In many PLs: `/` is both REAL and INTEGER division.
- If both arguments are INTEGER, it is INTEGER division with INTEGER result.
- Assume SUM and COUNT are INTEGER, and AVG is REAL.
- `AVG = SUM / COUNT`
- `SUM / COUNT` is computed, result is **truncated** to INTEGER. Then it is assigned to AVG as a REAL value.
- Solution is to use a different symbol for integer division (e.g., `div`).

Overloaded Operators

- **Ada, C++ and FORTRAN 90 allow user defined operator overloading.**
- If $+$ and $*$ are overloaded for matrix data type:
 - $A * B + C * D$
 - can be written for
 - `MatrixAdd(MatrixMult(A, B), MatrixMult(C, D))`
- Potential problems:
 - Users can define nonsense operations
 - Readability may suffer, even when the operators make sense

Type Conversions

- A *narrowing conversion* is one that converts an object to a type that cannot include all of the values of the original type e.g., `float` to `int`
- A *widening conversion* is one in which an object is converted to a type that can include at least approximations to all of the values of the original type e.g., `int` to `float`

Type Conversions – Mixed mode expression

- A *mixed-mode expression* is one that has operands of different types
- **Coercion:** an implicit type conversion
- **Cast:** Explicit type conversion requested by the programmer.
- Type conversions (coercion or cast) are either narrowing or widening.
- **Narrowing:** convert into a subset
- (e.g., `double` to `float`, `float` to `int`)
- **Widening:** convert into a superset
- (e.g., `float` to `double`, `int` to `float`)

Type Conversions – Coercions

- In **FORTRAN77** all coercions are widening.
- For example, if in an expression operands are INTEGER and REAL, then the compiler converts INTEGER to REAL type.
- Note that FORTRAN77 does not require the compilers to type check the parameters of user defined functions.
- Example,
- INTEGER A, C
- FUN (I) = 2 * I *I is integer*
- A = 2
- D = 3.6
- C = FUN (A + D) *A+D = 5.6 is Real*
- PRINT *, C
- END
- Prints 10
- In the function call to FUN, A is converted to REAL, then A + D is evaluated as REAL.
- Since FUN takes INTEGER argument (I), it is truncated to INTEGER.
- Compiler does not indicate an error.

Type Conversions – coercions

- Disadvantage of coercions:
 - They decrease in the type error detection ability of the compiler
- In most languages, all numeric types are coerced in expressions, using widening conversions
- In Ada, there are virtually no coercions in expressions
- Ada, and Modula-2 do not allow integer and floating-point operands in an expression.
- Exception in Ada: `**` (exponentiation operator) can take float or integer as its first argument. The second argument is always integer.

Explicit Type Conversions

- Explicit Type Conversions
- Called *casting* in C-based language
- Both **Modula-2** and **Ada** provide explicit type conversions in the form of **function calls**.
- $AVG := \text{FLOAT}(\text{SUM}) / \text{FLOAT}(\text{COUNT})$
- Here, SUM and COUNT can be any numerical type.
- AVG is FLOAT.

Note that Ada's syntax is similar to function calls

- In **C**,
- $avg = (\text{float}) \text{ sum} / (\text{float}) \text{ count};$
- In **C++**, both the syntax of Ada and C are acceptable.

Type Conversions: Errors in Expressions

- Causes
 - Inherent limitations of arithmetic division by zero e.g.,
 - Limitations of computer arithmetic overflow e.g.
- Often ignored by the run-time system

Relational and Boolean Expressions

- **Relational Operator:** Compares the values of its operands
- **Relational Expression:** Two operands and a relational operator
- The value of a relational expression is **boolean**.
- Operator symbols used vary somewhat among languages
(`!=`, `/=`, `.NE.`, `<>`, `#`)

Operation	Pascal	Ada	C	FORTRAN
Equal	<code>=</code>	<code>=</code>	<code>==</code>	<code>.EQ.</code>
not equal	<code><></code>	<code>/=</code>	<code>!=</code>	<code>.NE.</code>
greater than	<code>></code>	<code>></code>	<code>></code>	<code>.GT.</code>
less than	<code><</code>	<code><</code>	<code><</code>	<code>.LT.</code>
greater than or equal	<code>>=</code>	<code>>=</code>	<code>>=</code>	<code>.GE.</code>

Relational and Boolean Expressions

- Boolean Expressions
 - Operands are Boolean and the result is Boolean
- Boolean expressions consist of
 - Boolean variables
 - Boolean constants
 - Relational expressions
 - Boolean operators (AND, OR, NOT)

- Example operators

FORTRAN 77	FORTRAN 99	C	Ada
.AND.	and	&&	and
.OR.	or		or
.NOT.	not	!	not
			xor

Relational and Boolean Expressions

- C has no Boolean type--it uses `int` type with 0 for false and nonzero for true
 - One odd characteristic of C's expressions:
 - **`a < b < c`** is a legal expression, but the result is not what you might expect:
 - In C, relational operators are **left associative**.
 - `a > b > c`
 - means evaluate first **`a > b`**, resulting 0 or 1, then compare this result (0, or 1) with c.
 - that is the result is `1 > c` or `0 > c`, depending on `a > b`.
 - Whereas, commonsense interpretation is that `a > b` and `b > c`.
 - `main()` {
 - `printf (“%d\n”, 6 > 4 > 1);`
 - `}` *prints 0.*
 - **Readability** requires a PL to include **boolean type**.
-

Relational and Boolean Expressions: Operator Precedence

- **Precedence:** NOT (highest) AND OR (lowest)
- Arithmetic, relational and boolean operators can all be in the same expression.
- A PL must define the precedence of all operators.

Relational and Boolean Expressions: Operator Precedence

- Precedence of C-based operators

prefix ++, --

unary +, -, prefix ++, --, !

*, /, %

binary +, -

<, >, <=, >=

=, !=

&&

||

Relational and Boolean Expressions: Operator Precedence

In **FORTRAN**:

Arithmetic (highest)

Relational

Boolean (lowest)

Example,

$A + B .GT. 2 * C .AND. K .NE. 0$

is evaluated as

$(2 * C)$

$(A + B)$

$((A + B) .GT. (2 * C)) (K .NE. 0)$

$((A + B) .GT. (2 * C)) .AND. (K .NE. 0)$

Relational and Boolean Expressions: Operator Precedence

Example:

```
LOGICAL EXPR
```

```
A=2
```

```
B=3
```

```
C=4
```

```
K=5
```

```
EXPR = A + B .GT. 2 * C .AND. K .NE. 0
```

```
PRINT *, EXPR
```

```
A=10
```

```
EXPR = A + B .GT. 2 * C .AND. K .NE. 0
```

```
PRINT *, EXPR
```

```
END
```

C Output is

C F

C T

Boolean constants in FORTRAN are .TRUE. and .FALSE.

Short Circuit Evaluation

- An expression in which the result is determined without evaluating all of the operands and/or operators
- Example: $(13 * a) * (b / 13 - 1)$
If a is zero, there is no need to evaluate $(b / 13 - 1)$
- Problem with non-short-circuit evaluation


```
index = 1;
while (index <= length) && (LIST[index] != value)
    index++;
```

 - When $\text{index} = \text{length}$, $\text{LIST}[\text{index}]$ will cause an indexing problem (assuming LIST has $\text{length} - 1$ elements)

Short Circuit Evaluation

- C, C++, and Java: use short-circuit evaluation for the usual Boolean operators (**&&** and **||**), but also provide bitwise Boolean operators that are not short circuit (**&** and **|**)
- Most standard Pascal compilers do not use short-circuit evaluation, instead they evaluate all operands.
- Short-circuit evaluation exposes the potential problem of side effects in expressions
e.g. `(a > b) || (b++ / 3)`
- Here, `b` is incremented only if `a <= b`, If the programmer assumes that `b` is incremented every time, this is an error.
- In **Ada**, user can define short-circuit by the `and then` and `or else` **operators**.
- `I := 1;`
- **while** `(I <= LISTLEN) and then (LIST[I] /= key)`
- `loop`
- `I := I + 1`
- **end loop**
- In **C** and **Modula-2** every evaluation of AND and OR is **short-circuit**.

Assignment Statements

- The general syntax
 <target_var> <assign_operator> <expression>
- The assignment operator
 - = FORTRAN, BASIC, PL/I, C, C++, Java
 - : = ALGOLs, Pascal, Ada
- = can be bad when it is overloaded for the relational operator for equality
- In **PL/I**, the symbol = is both **assignment** and **relational** operator.
- A = B = C
- assigns A the value TRUE if B = C, FALSE otherwise.

Assignment Statements: Multiple Targets

- In **PL/I**, the statement
- $A, B = 0$
- assigns the value 0 to both A and B.

Assignment Statements: Conditional Targets

- Conditional targets (C, C++, and Java)
`(flag ? count1 : count2) = 0`

Which is equivalent to

```
if (flag)
  count1 = 0
else
  count2 = 0
```

- Be careful, without the parantheses
- `flag ? count1 : count2 = 0`
- is equivalent to
- `if flag then count1`
- `else count2 = 0`

Assignment Statements: Compound Operators

- A shorthand method of specifying a commonly needed form of assignment
- Introduced in ALGOL; adopted by C
- Example

$a = a + b$

is written as

$a += b$

Assignment Statements: Unary Assignment Operators

- Used in C,
- `count++;` is equivalent to `count = count + 1;`
- Unary assignment operators in C-based languages combine increment and decrement operations with assignment
- `sum = ++count;` is equivalent to
 - `count = count + 1;`
 - `sum = count;`
- `sum = count++;` is equivalent to
 - `sum = count;`
 - `count = count + 1;`

Assignment Statements: Unary Assignment Operators

- When two unary operators apply to the same operand, the association is from **right to left**.
- `-count++` is equivalent to
- `count = count+1; -count;`
- Unary operations may cause unreadability.

Assignment Statements: Unary Assignment Operators

Consider the following program

```
main()
{ int b=5;
  b = b+++b++; /* legal */
  printf(“%d\n”, b);
}
```

- The behavior of such a program is not defined clearly in the C language. So, different compilers give different results.
- For example, when compiled with cc it gives **12**, whereas
- when compiled with gcc it gives **6**.

Assignment as an Expression

- In C, C++, and Java, the assignment statement produces a result and can be used as operands
- An example:

```
while ((ch = getchar()) != EOF) {...}
```

`ch = getchar()` is carried out; the result (assigned to `ch`) is used as a conditional value for the `while` statement

Assignment as an Expression

- In **C**, **assignment** statement **produces a value** as if an operator.
- $a = b = c$
- a and b get the value of c .
- Good for initializing a set of variables to the same value.
- Difficulty:
 - `if (x=y)`
 - is true if $y > 0$, and assigns y to x , but the programmer meant
 - `if (x==y)`
- Compiler cannot detect such common typing errors.

Mixed mode assignments

- Assignment statements can also be mixed-mode, for example

```
int a, b;
```

```
float c;
```

```
c = a / b;
```

- In Pascal, integer variables can be assigned to real variables, but real variables cannot be assigned to integers
- In Java, only widening assignment coercions are done
- In Ada, there is no assignment coercion