# Statement level control structures

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#### Control Statements: Evolution

- FORTRAN I control statements were based directly on IBM 704 hardware
- Much research and argument in the 1960s about the issue
  - One important result: It was proven that all algorithms represented by flowcharts can be coded with only twoway selection and pretest logical loops

#### Control structure

#### • Control Structure:

- a control statement
- And the collection of statements whose execution it controls.
- Control Statements are:
- **Conditional** execution statements
- **Repeated** execution statements
- Control structures should have single entry and single exit points.

## Compound statements

- ALGOL60 introduced the first statement collection structure.
- First introduced in ALGOL 60 in the form of
- begin
- statement<sub>1</sub>
- •
- statement
- end
- A collection of statements is treated **as a single statement**.
- ALGOL60 allows **data declarations** at the beginning of a compound statement making it a **block**.
- begin
- integer v1, v2;
- •
- end

. . .

• Scope of v1 and v2 is the block.

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- A *selection statement* provides the means of choosing between two or more paths of execution
- Two general categories:
  - Two-way selectors
  - Multiple-way selectors

- General form: if control\_expression then clause else clause
  - Design Issues:
    - What is the form and type of the control expression?
    - How are the **then** and **else** clauses specified?
    - How should the meaning of nested selectors be specified?

- A special case: **Single-way selector**: In most cases a subform of a two-way selector. FORTRAN IV has no two-way selector.
- In FORTRAN this is called logical IF
- FORTRAN: **IF** (boolean\_expr) statement
- Problem: can select only a single statement; to select more, a GOTO must be used, as in the following example
  - IF (.NOT. condition) GOTO 20

20 CONTINUE

- Negative logic is bad for readability
- This problem was solved in FORTRAN 77

#### Selection Statements

- Example,
- IF (A .GT. B) GOTO 10
- A = B
- C = A + D
- 10 CONTINUE
- •
- IF (A .GT. 0) A = 0

- Most later languages allow compounds for the selectable segment of their single-way selectors
- The concept of **compound statements** is introduced by **ALGOL 60**.
- ALGOL60 allows compound statements in selection.
- if (Boolean Expression) then
- begin
- statement<sub>1</sub>
- •
- statement<sub>n</sub>
- end

#### Two-Way Selection: Examples

- ALGOL60 introduced the first two-way selector.
- if (boolean\_expr)

**then** statement (then clause)

**else** statement (else clause)

• The statements could be single or compound

Java example

if (sum == 0)
if (count == 0)
result = 0;

else result = 1;
Which if gets the else? AMBIGUOUS

• To force an alternative semantics, compound statements may be used:

```
if (sum == 0) {
```

```
if (count == 0)
```

```
result = 0;
```

else result = 1;

#### Nesting selectors

- PASCAL, C, C++, Java: match else to the nearest unmatched if.
   Only a semantic rule, not a syntactic rule (not clear from the
  - syntax).
- Perl requires all then and else statements to be compound.
- In ALGOL60, an if statement is **not allowed** to be **nested directly in** a **then** clause. So, *the code above is illegal*.
  - The programmer must use compound statements.

#### **Special Words and Selection Closure**

- Another solution to the problem above is to use **special words**.
- **IF** marks the beginning.
- **THEN** marks the end of condition and the beginning of then part.
- ELSE marks the end of then part and the beginning of else part.
- END or END IF marks the end of IF.
- In Modula-2, If-then-else construct has an END, even if a single stmt.
- If cond If cond
- THEN THEN
- ts1 ts1
- ... ... ... • tsn tsn
- ELSE END
- es1
- ...
- Esm
- END
- In FORTRAN77, FORTRAN90 and Ada, the closing of an IF is END IF.
- This is even more **readable** than Modula-2.

- Allow the selection of one of any number of statements or statement groups
- Design Issues:
  - 1. What is the form and type of the control expression?
  - 2. How are the selectable segments specified?
  - 3. Is execution flow through the structure restricted to include just a single selectable segment?
  - 4. What is done about unrepresented expression values?



- FORTRAN's computed GOTO
- GOTO (label, label, ..., label) integer\_expression
- Evaluate *integer expression*. If it is  $\overline{i}$ , go to the statement with label label.
- If no label for the value of the expression is given go to the next statement.

- Modern Multiple Selectors
- case statement in ALGOL-W,
- **case** integer\_expression **of**
- begin
- *statement*<sub>1</sub>;
- .
- *statement*<sub>n</sub>;
- end
- If integer\_expression is *i*, evaluate statement<sub>i</sub>.

- case statement in PASCAL (selectable statements are labeled)
- case ordinal\_type\_expression of
- constant\_list\_1: statement\_1;
- *constant list n: statement n;*
- end
- A constant may not appear in more than one constant list.
- Undefined results occur if the value of the expression is not in any of the lists (standard Pascal).
- ANSI/IEEE Pascal Standard specifies that the code generated should detect such cases and **report error messages**.
- Many dialects of PASCAL now include **else** to match any **unlisted** value.

```
Example:
Case idx of
1,3: a := a+1;
2: begin
b := b+1;
c := c-1;
end;
else
writeln ("idx: ", idx, " is strange.");
end;
```

```
switch in C: relatively primitive
switch (expression){/* expression is integer type */
case constant_expression_1: statement_1;
...
```

```
case constant_expression_n: statement_n;
[default: statement_n+1;]
```

Control of execution is transferred to the statement whose *constant expression* is equal to the *expression*. Then all following cases are executed. **break** statement is used to avoid the execution of unwanted cases.

#### Design choices for C's switch statement

- **1.** Control expression can be only an integer type
- 2. Selectable segments can be statement sequences, blocks, or compound statements
- **3.** Any number of segments can be executed in one execution of the construct (there is no implicit branch at the end of selectable segments)
- 4. default clause is for unrepresented values (if there is no default, the whole statement does nothing)

```
• The Ada case statement
case expression is
when choice list => stmt_sequence;
...
when choice list => stmt_sequence;
when others => stmt_sequence;]
end case;
```

More reliable than C's switch (once a stmt\_sequence execution is completed, control is passed to the first statement after the case statements.

• Multiple Selectors can appear as direct extensions to two-way selectors, using else-if clauses,

if ...
then ...
elsif ...
elsif ...
then ...
else ...
end if

- Ada provides a special case of **nested if** along with a **case** statement.
- **if** *boolean\_expression\_1* **then** *statement\_1*;
- elsif boolean\_expression\_2 then statement\_2;
- elsif boolean\_expression\_3 then statement\_3;

•

- end if;
- This is more readable than standard nested if and case structures.

- The repeated execution of a statement or compound statement is accomplished either by iteration or recursion
- General design issues for iteration control statements:
- 1. How is iteration controlled? logical counting a combination of the two
- 2. Where is the control mechanism in the loop?pretest: at the top orposttest: at the bottom

- Control statement has a variable called **loop variable**, along with its **initial** and **terminal** values, and sometimes a **stepsize**, they are called **loop parameters**.
  - Design Issues
  - type and scope of the loop variable?
  - value of the loop variable at the loop termination?
  - can the **loop variable** or loop parameters be **changed** in the loop?
  - pretest or posttest?
  - loop parameters are evaluated only once or for every iteration?

```
The DO Statement of FORTRAN77 and FORTRAN 90
```

**DO** *label variable* = *initial*, *terminal* [, *stepsize*]

**Loop variable** can be **integer**, **real** or **double-precision**. DO 10 R=0.5,9.9,0.1

Iteration count is computed before the execution.

- Parameters can be changed in the loop, but iteration count remains unchanged.
- Gnu implantation of FORTRAN 77 does not allow the loop counter to be modified in the loop. Following code does not compile in g77. Do 10 J=1,10

$$\mathbf{J}=\mathbf{3}$$

#### **10 CONTINUE**

It is compiled by f77 (SUN compiler), prints infinite 5's (infinite loop). IN FORTRAN space character is not used as a token separator. That is, DO 10 J=1, 10, and DO10J=1, 10 are the same. If you type "." for ", " it is DO10J=1.10, an assignment.

```
FORTRAN 90 DO includes
```

[name] DO variable = initial, terminal [, stepsize]

```
END DO [name]
Fortran DO statement is pretest.
The following program produces no output.
 DO 20 I =4,3
 Print *, I
20 continue
end
However, the following program
 DO 20 I =4, 3, -1
20 Print *, I
end
Produces the following output:
4
3
```

The ALGOL60 for Statement

A significant generalization of FORTRAN's DO: User can **combine** a **counter** and a **boolean expression** for the loop control. However, **Flexibility** leads to **complexity**.

#### Syntax in EBNF:

```
<for_stmt>-> for var := <list_elt>{,<list_elt>} do <statement>
<list_elt> -> <expression>
| <expression> step <expression> until <expression>
| <expression> while <Boolean expression>
```

**Combines** a **counter** and a **Boolean expression** for loop control. Example, the following statements are equivalent.

```
for i := 1,2,3,4,5 do list[i] := 0
for i := 1 step 1 until 5 do list[i] := 0
for i := 1, i+1 while (i <= 5) do list[i] := 0</pre>
```

Much more complex, yet more flexible, then any other for loop.

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```
Code for step-until
for var <- init expr
loop:
form
  Until <- until exp
  step <- step exp</pre>
  tmp <- (for var - until) * SIGN(step)</pre>
  if tmp > 0 go to out
  [loop body]
  for var <- for var + step
  go to loop
out: ...
```

- Example,
- for i := 1, 2, 5 step 2 until 11, 2\*i while i<90, 41, -5 do
- will execute for the following values of i:
- 1, 2, 5, 7, 9, 11, 26, 52, 41, -5
- Too difficult to understand
- All expressions are evaluated for every iteration
- pretest loop
- Loop variables can be integer or real
- Loop variables can not be changed in the loop body

```
<for_stmt>-> for var := <list_elt>{,<list_elt>} do <statement>
<list_elt> -> <expression>
| <expression> step <expression> until <expression>
| <expression> while <Boolean_expression>
```

#### for i := 1, 2, 5 step 2 until 11, 2\*i while i<90, 41, -5 do

will execute for the following values of i:

```
1, 2, 5, 7, 9, 11, 26, 52, 41, -5
```

```
Example,
i := 1;
for c := 1 step c until 3*i do i:=i+1
```

#### will execute as

The Pascal for Statement

The model of simplicity for <var> := <init\_exp> (to | down to) <final\_exp> do <stmt>

The Ada for Statement Relatively **simple**, **pretest** loop.

**for** *variable* **in** [**reverse**] *discrete\_range* **loop** ... end loop

Scope of the variable is the range of the loop.

```
SUM : FLOAT := 0;COUNT : FLOAT := 1.35;for COUNT in 1..10 loopSUM := SUM + COUNTthis COUNT is INTEGERend loopFLOAT COUNT is hidden here
```

here COUNT is 1.35, SUM is 55

• C's for statement

for ([expr\_1] ; [expr\_2] ; [expr\_3]) statement

- The expressions can be whole statements, or even statement sequences, with the statements separated by commas
  - The value of a multiple-statement expression is the value of the last statement in the expression
- There is no explicit loop variable
- Everything can be changed in the loop
- The first expression is evaluated once, but the other two are evaluated with each iteration

The for Statement of C, C++, and Java

```
• pretest
```

for (exp1; exp2; exp3) statement

Statement can be single, compound or null. exp1 is evaluated only once when execution begins. exp2 is evaluated before each execution of the loop. loop terminates when exp2 is 0

**exp3** is evaluated **after each** execution of the loop.

```
exp1
loop:
  if exp2 = 0 go to out
  [loop body]
  exp3
  go to loop
out:...
```

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- All of the expressions (exp's) of C's for are optional.
- So, for (;;) stmt; is legal.
- An absent exp2 is considered true.
- No explicit loop variable or parameters
- Variables can be changed in the body
- Each expression can comprise **multiple statements** separated by , the value of the expression is the value of the last statement.

- C++ differs from C in two ways:
  - 1. The control expression can also be Boolean
  - 2. The initial expression can include variable definitions (scope is from the definition to the end of the loop body)
- Java and C#
  - Differs from C++ in that the control expression must be Boolean

#### Iterative Statements: Logically-Controlled Loops

- More general than counter-controlled loops
- Repetition control is based on a Boolean
- Design issues:
  - Pre-test or post-test?
  - Should the logically controlled loop be a special case of the counting loop statement ? expression rather than a counter
- General forms:

```
while (ctrl_expr) do
loop body loop body
while (ctrl_expr)
```

#### Iterative Statements: Logically-Controlled Loops

- Pascal has separate pre-test and post-test logical loop statements (while-do and repeat-until)
  - Pascal **repeat-until** (posttest) loop can have a single statement, compound statement, or statement sequence. This is the only control structure with this flexibility. This another reason for **lack of orthogonality** in Pascal.
- C and C++ also have both, but the control expression for the post-test version is treated just like in the pre-test case (while-do and do- while)
- Java is like C, except the control expression must be Boolean (and the body can only be entered at the beginning -- Java has no **goto**

#### Iterative Statements: Logically-Controlled Loops

- Ada has a pretest version, but no post-test
- FORTRAN 77 and 90 have neither
- Perl has two pre-test logical loops, while and until, but no post-test logical loop

# Iterative Statements: User-located loop control

- Sometimes it is convenient for the programmers to decide a location for loop control (other than top or bottom of the loop)
- Simple design for single loops (e.g., break)
- Design issues for nested loops
  - 1. Should the conditional be part of the exit?
  - 2. Should control be transferable out of more than one loop?

#### Iterative Statements: User-Located Loop Control Mechanisms break and continue

- C, C++, and Java: **break** statement
- Unconditional; for any loop or switch; one level only
- Java and C# have a labeled break statement: control transfers to the label
- An alternative: **continue** statement; it skips the remainder of this iteration, but does not exit the loop

## Iterative Statements: User-Located Loop

Both FORTRAN90 and Ada have loops with no control (infinite loop). Ada example,

loop

... end loop

and an **exit** statement is **provided**. **exit** [*loop-label*] [**when** *condition*]

# Iterative Statements: User-Located Loop

- In Ada, loops can have labels.
- If *loop-label* is omitted, exit statement **causes** the **termination** of **only** the **block** in which **it appears**.
- Transfers control to the statement after the loop.

OL: loop

```
...
IL: for ... loop
...
exit OL when x>0;
...
end loop IL;
...
end loop OL;
```

(exits the OL loop)

# Iterative Statements: User-Located Loop

C and FORTRAN90 have statements to skip the rest of a single iteration. In C, continue, in FORTRAN 90 CYCLE. while ( ... ) {



Such statements **make** a language **unreadble**.

# Iterative Statements: Iteration Based on Data

- Number of elements of in a data structure control loop iteration
- Control mechanism is a call to an *iterator* function that returns the next element in some chosen order, if there is one; else loop is terminate
- C's for can be used to build a user-defined iterator: for (p=root; p==NULL; traverse(p)) { }

# Iterative Statements: Iteration Based on Data

• C#'s foreach statement iterates on the elements of arrays and other collections:

```
Strings[] = strList = {"Bob", "Carol", "Ted"};
foreach (Strings name in strList)
Console.WriteLine ("Name: {0}", name);
```

• The notation  $\{0\}$  indicates the position in the string to be displayed

#### **Unconditional Branching**

- Transfers execution control to a specified place in the program
- Represented one of the most heated debates in 1960's and 1970's
- Well-known mechanism: goto statement
- Major concern: Readability
- Some languages do not support goto statement (e.g., Module-2 and Java)
- C# offers goto statement (can be used in switch statements)
- Loop exit statements are restricted and somewhat camouflaged goto's

- Suggested by Dijkstra
- Purpose: to support a new programming methodology that supported verification (correctness) during development
- Basis for two linguistic mechanisms for concurrent programming (in CSP and Ada)
- Basic Idea: if the order of evaluation is not important, the program should not specify one

• Form

- if <Boolean exp> -> <statement>
- [] <Boolean exp> -> <statement>
- ...
  [] <Boolean exp> -> <statement>
  fi
- Semantics: when construct is reached,
  - Evaluate all Boolean expressions
  - If more than one are true, choose one non-deterministically
  - If none are true, it is a runtime error

#### Selection Guarded Command: Illustrated



#### Selection Guarded Command

- To find the largest of the two values:
- if  $x \ge y \ge max := x$
- []  $y \ge x -> max := y$
- fi
- A good way to state that the order of execution is irrelevant.

#### Loop Guarded Command

#### • Form

do <Boolean> -> <statement>

[] <Boolean> -> <statement>

```
...
[] <Boolean> -> <statement>
od
```

- Semantics: for each iteration
  - Evaluate all Boolean expressions
  - If more than one are true, choose one nondeterministically; then start loop again
  - If none are true, exit loop

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#### Loop Guarded Command: II



#### Guarded Commands: Rationale

- Connection between control statements and program verification is intimate
- Verification is impossible with goto statements
- Verification is possible with only selection and logical pretest loops
- Verification is relatively simple with only guarded commands