

# Systems Programming

Chapter 1  
Background

## Examples of System Software

- text editor: create/modify a program in a high level language
- compiler: translate it to machine code
- loader/linker: load resulting code into memory and prepare for execution
- debugger: detect errors in program

## Outline

- Introduction to system software
- An overview of the material
- Relationship bw system software and machine architecture
- Simplified Instructional Computer (SIC)
- Architecture of several computers

## Examples of System Software

- text editor: create/modify a program in assembly language
  - w/macro instructions to read and write data, or to perform other higher-level functions
- assembler/macro processor: translate it to machine code
- loader/linker: load resulting code into memory and prepare for execution
- debugger: detect errors in program

## System Software

- a variety of programs that support the operation of a computer
- makes it possible for the user to focus on an application or other problem to be solved, w/o needing to know the details of how the machine works internally.

## Examples of System Software

- Operating system
  - UNIX, DOS: textual user interface
  - MacOS, Windows: GUI (graphical user interface)
  - takes care of all machine-level details for us
    - connect to network, use different kinds of devices, perform input/output, etc.

## We will . . .

- understand the processes that were going on “behind the scenes” as you used the computer
- gain a deeper understanding of how computers actually work

## Machine-Independent Aspects

- Assembler
  - general design and logic
- Compilers
  - some code optimization techniques
- Linkers
  - linking independently assembled subprograms does not usually depend on the computer being used

## System Software and Machine Architecture

### System software

- machine dependent
- intended to support the operation and use of the computer itself

### Application software

- machine independent
- intended to support a particular application

## Simplified Instructional Computer (SIC)

- a hypothetical computer that has been carefully designed to include the hardware features most often found in real machines
  - unusual or irrelevant complexities of real machines have been avoided

## Machine-Dependent Aspects

- Assemblers:
  - translate mnemonic instructions into machine code
    - the instruction formats, addressing modes, etc. are of direct concern in assembler design
- Compilers
  - generate machine language code
    - number and type of registers and machine instructions available
- Operating Systems
  - manage the resources of a computer

## Organization of Chapters

- Fundamental features of system software being discussed
- Machine-dependent features
- Machine-independent features
- Major design options for structuring a particular piece of software
- Examples of implementations on actual machines

## Simplified Instructional Computer (SIC)

- Two versions
  - standard model
  - XE model
- They are upward-compatible: An object program for the standard SIC machine will also execute properly on a SIC/XE system.

- Memory:
  - $2^{15}$  bytes, 3-byte words, byte addresses
- Registers:
  - 5 3-byte long registers

Mnemonic	Number	Special Use
A	0	Accumulator; used for arithmetic operations
X	1	Index register; used for addressing
L	2	Linkage register; the Jump to Subroutine (JSUB) instruction stores the return address
PC	8	Program Counter; the address of next instruction to be fetched
SW	9	Status Word; contains a variety of information, including a Condition Code (CC)

## QUESTION

- How can you characterize a machine architecture?

- Data formats:
  - 24-bit Integers (2's complement form for negatives)
  - 8-bit ASCII chars
- Instruction formats:
  - 24-bit instructions



- Addressing modes:
- | Mode    | Indication | Target Address Calculation |
|---------|------------|----------------------------|
| Direct  | x = 0      | TA = address               |
| Indexed | x = 1      | TA = address + (X)         |

## SIC Machine Architecture

- Memory
- Registers
- Data formats
- Instruction formats
- Addressing modes
- Instruction set
- Input and Output

- Instruction Set:
  - load and store registers: LDA, LDX, STA, STX, etc.
  - Integer arithmetic operations: ADD, SUB, MUL, DIV
    - all involve register A and a word in memory, with the result left in the register
  - COMP: compare the value in A with a word in memory, set the condition code (CC) to indicate the result (<, =, >)
  - Conditional jump instructions: JLT, JEQ, JGT
  - Subroutine linkage:
    - JSUB jumps to subroutine, places return address in register L
    - RSUB returns by jumping to the address contained in register L

• Input and Output:

- performed by transferring 1 byte at a time to or from the rightmost 8 bits of register A
- Each device is assigned a unique 8-bit code.
- Test Device (TD) instruction tests whether the addressed device is ready to send or receive a byte of data and sets CC (< ready, = not ready)
- Read Data (RD) reads a byte
- Write Data (WD) writes a byte

### SIC Sample Arithmetic Operations

- All arithmetic operations are performed using register A, with the result being left in register A.

```

LDA  ALPHA  LOAD ALPHA INTO REGISTER A
ADD  INCR   INCREMENT THE VALUE IN REGISTER A
SUB  ONE    SUBTRACT ONE
STA  BETA   STORE IN BETA
LDA  GAMMA  LOAD GAMMA INTO REGISTER A
ADD  INCR   ADD THE VALUE OF INCREMENT
SUB  ONE    SUBTRACT 1
STA  DELTA  STORE IN DELTA      :
:
:
ONE  WORD  1      ONE-WORD CONSTANT
ALPHA RESW  1     ONE-WORD VARIABLES
BETA  RESW  1
GAMMA RESW  1
DELTA RESW  1
INCR  RESW  1
    
```

### SIC/XE Machine Architecture

- More Memory
- Extra Registers
- Data formats – floating point numbers
- Additional Instruction formats
- More Addressing modes
- Extended Instruction set
- Input and Output – channels to overlap I/O and processing

### SIC Looping and Indexing Operations

- X is the index register.

```

LDX  ZERO   INITIALIZE INDEX REGISTER TO 0
MOVECH LDCH STR1,X LOAD CHARACTER FROM STR1 INTO REG A
STCH  STR2,X STORE CHARACTER INTO STR2
TIIX  ELEVEN ADD 1 TO INDEX, COMPARE RESULT TO 11
JLT   MOVECH LOOP IF INDEX IS LESS THAN 11
ADD   INCR   ADD THE VALUE OF INCREMENT
:
:
:
STR1  BYTE  C'TEST STRING'  11-BYTE STRING CONSTANT
STR2  RESB  11              11-BYTE VARIABLE
:
:
ZERO  WORD  0
ELEVEN WORD 11
    
```

### SIC Sample Data Movement Operations

- there are no memory-to-memory move instructions

```

LDA  FIVE   LOAD CONSTANT 5 INTO REGISTER A
STA  ALPHA  STORE IN ALPHA
LDCH CHARZ  LOAD CHARACTER 'Z' INTO REGISTER A
STCH C1     STORE IN CHARACTER VARIABLE C1
:
:
:
ALPHA RESW  1      ONE-WORD VARIABLE
FIVE  WORD  5      ONE-WORD CONSTANT
CHARZ BYTE  C'Z'   ONE-BYTE CONSTANT
C1    RESB  1      ONE-BYTE VARIABLE
    
```

Four ways of defining storage

### SIC Looping and Indexing Operations 2

- Add together corresponding elements of ALPHA and BETA, store results in GAMMA.

```

LDA  ZERO   INITIALIZE INDEX VALUE TO 0
STA  INDEX
ADDLP LDX   INDEX  LOAD INDEX VALUE INTO REG X
LDA  ALPHA,X LOAD WORD FROM ALPHA INTO REG A
ADD  BETA,X  ADD WORD FROM BETA
STA  GAMMA,X STORE THE RESULT IN A WORD IN GAMMA
LDA  INDEX  ADD 3 TO INDEX VALUE
ADD  THREE
STA  INDEX
COMP K300   COMPARE NEW INDEX VALUE TO 300
JLT  ADDLP  LOOP IF INDEX IS LESS THAN 300
:
:
INDEX RESW  1      ONE-WORD VARIABLE FOR INDEX VALUE
                    ARRAY VARIABLES - 100 WORDS EACH
ALPHA RESW  100   : ONE-WORD CONSTANTS
BETA  RESW  100   ZERO WORD 0
GAMMA RESW  100   K300 WORD 300
                    THREE WORD 3
    
```

### SIC Input/Output Operations

- Read 1 byte of data from device F1 and copy it to device 05.

```

INLOOP TD      INDEV  TEST INPUT DEVICE
JEQ  INLOOP   LOOP UNTIL DEVICE IS READY
RD    INDEV   READ ONE BYTE INTO REGISTER A
STCH  DATA   STORE BYTE THAT WAS READ
:
OUTLP TD      OUTDEV  TEST OUTPUT DEVICE
JEQ  OUTLP   LOOP UNTIL DEVICE IS READY
LDCH  DATA   LOAD DATA BYTE INTO REGISTER A
WD    OUTDEV  WRITE ONE BYTE TO OUTPUT DEVICE
:
INDEV BYTE    X'F1'  INPUT DEVICE NUMBER
OUTDEV BYTE   X'05   OUTPUT DEVICE NUMBER
DATA  RESB   1      ONE-BYTE VARIABLE
    
```

### RISC (Reduced Instruction Set Computers)

- Developed in early 1980s to simplify the design of processors
  - faster, less expensive processors
  - greater reliability
  - faster instruction execution times
- Fixed instruction length (usually 1 word)
  - Single cycle execution for most instructions
- All instructions, except Load and Store, are register-to-register
  - Memory access by Load and Store instrs only
  - Have large number of general purpose registers
- Number of machine instructions, instruction formats, and addressing modes are small.
- Examples: UltraSPARC, PowerPC.

### SIC Subroutine Call and Record Input Operations

- Read a 100-byte record from an input device into memory.

```

JSUB  READ    CALL READ ROUTINE
:
READ  LDX     ZERO   SUBROUTINE TO READ 100 -BYTE RECORD
RLOOP TD     INDEV  TEST INPUT DEVICE
JEQ  RLOOP   LOOP UNTIL DEVICE IS READY
RD    INDEV  READ ONE BYTE INTO REGISTER A
STCH  RECORD,X STORE DATA BYTE INTO RECORD
TIK  K100   ADD 1 TO INDEX AND COMPARE TO 100
JLT  RLOOP   LOOP IF INDEX IS LESS THAN 100
RSUB  :      EXIT FROM SUBROUTINE
:
INDEV BYTE    X'F1'  INPUT DEVICE NUMBER
RECORD RESB  100    100-BYTE BUFFER FOR INPUT RECORD
:
ZERO  WORD    0
K100  WORD    100   ONE-WORD CONSTANTS
    
```

### Main Machine Architectures

- CISC (Complex Instruction Set Computers)
  - relatively large and complicated instruction set
  - several different instruction formats and lengths
  - many different addressing modes
  - requires complicated hardware mechanisms
  - Examples: VAX, Pentium Pro
- RISC (Reduced Instruction Set Computers)